


## saia'p



- Four standard versions with up to 64 I +0
- 4K user steps
- Programmable according to ladder diagram, logic diagram, flow-chart or functional diagram as per DIN-standard.
- Arithmetic instructions and security commands for permanent monitoring of the operation (watchdog and check sum).
- Easy beginning for learners but with big power reserve for the demanding PC-user.
$v$
$\smile$
v
$\checkmark$


## saia’p

## Who uses the PCAD-manual?

We do not know how familiar you are with programmable controllers. Maybe you are a beginner or already an experienced PC-specialist.

This manual serves as a course, in order to give the beginner an easy introduction to the world of programmable controllers. Read especially the chapters 1 to 5 carefully before you start to work and do not let yourself be confused with the diagrams in the chapters 6 and 7. These are not important at the beginning if you have equipped yourself with the simulation and power supply unit PCA2.Sø5 (see chapter 5 E ).

In chapter 9 you will be gradually led up the staircase from A via B to C. For this, we use easy and clear examples, which can be collectively tested on your desk with the above-mentioned simulation unit.

If you do not know how to go on, please make use of the experience of our specialist in your vicinity or register for the next workshop.

We wish you a lot of fun with the versatile PCAØ!
If you are a PC-specialist, you do not have to read all the information in chapter 9. In this case, concentrate on the instruction lists at the beginning of parts $A, B$ and $C$. If you want to know more about these instructions, refer to the elaborated "Basic Manual".


1. A look at the exterior and interior of the PCAD

The PCAD is the compact series of the SAIA-PC system family. The input and output assemblies are combined on a single pc-board. Owing to the high intelligence very simple as well as quite complex problems can be solved with the PCAD.

On the exterior, the following functional parts are distinguished:


The following functional units can be seen inside:
ramming unit


The programming unit serves to transfer the user program to the user memory (RAM). The CPU executes this program, interrogates the states of the process inputs and controls the process outputs accordingly.
2. Common technical data

| Microprocessor system | 8ø85-2 |
| :---: | :---: |
| Cycle time per user instruction (average) | $7 \emptyset \mu s$ |
| Instruction set level (14) | 32 basic instructions + $2 \emptyset$ additional instructions incl. arithmetic |
| Number of parallel programs | 16 |
| Number of index registers | 16 (1 per parallel program) |
| Number of subroutine levels | 3 |
| User memory | 4K program lines ( $=8 \mathrm{~K}$ Bytes) |
| Volatile/non-volatile flags | 477* + $235=712$ |
| Number of software counters and timers | 64 addresses ( $C=64, T=32$ ) |
| Counting capacity | 65*535 |
| Time ranges (time base $\emptyset .1$ or $\emptyset .01 \mathrm{~s}$ ) | $\emptyset .1$ ( 0.01 ) to 650ø (650)s |
| Hardware timer PCAØ. Н2Ø | 4 time ranges $\emptyset .9 / 3.7 / 3 \emptyset / 24 \emptyset$ s |
| Connection of peripherals | via 25-pole PGU-connector |
| Operating modes | RUN, BREAK, STEP, MAN, PROG |
| Indicating lamps | LED for RUN/CPU RUN/WATCHDOG LED for I/O |
| $\begin{aligned} & \text { Inputs } \\ & (B 9 \emptyset) \end{aligned}$ | galvanically connected, source or sink operation nominal +24VDC $H=+19 \ldots+32 V$ <br> $L=\emptyset \ldots+4 V$ <br> $\mathrm{I}=10 \mathrm{~mA}, 24 \mathrm{VDC}, \mathrm{t}_{\mathrm{I}}=9 \mathrm{~ms}$ |
| Relay outputs (A21) | ```galvanically isolated, normally open contacts contact rating 3A, 250VAC AC1 3A, 24VDC DC1``` |
| Transistor outputs (B9 D $^{\text {) }}$ | galvanically connected, positive switching Ø.5mA. ..Ø.5A, 5...36VDC |
| Supply voltage | 24VDC $\pm 20 \%$ |
| Ambient temperature | $\emptyset \ldots+50^{\circ} \mathrm{C}$ |
| High noise immunity | as per IEC 255-4/E5 class III, i.e. 25øøV and IEC 8ø1-4 <br> class III (20ØØV) |

[^0]3. The four standard versions
3.1 With relay outputs

Type PCAØ.M12R M4
16 I, 24VDC
8 0, relay contact
max. 3A, 25ØVAC


### 3.2 With transistor outputs

Type PCAØ.M12T M4
$2 \emptyset$ I, 24VDC
120, $0.5 \mathrm{~A} / 24 \mathrm{VDC}$


Type PCAØ.M14R M4
32 I, 24VDC
160 , relay contact
max. 3A, 25ØVAC


Type PCAØ.M14T M4
$4 \emptyset$ I, 24VDC
250, D. 5A/24VDC


Detailed information on the inputs/outputs see chapter 7 .
Please note: If a certain number of units are ordered, we will supply you, too, with a custom-made version. Please contact our nearest selling agency.
4. Important accessories

### 4.1 The hardware timer module PCAØ. $\mathrm{H} 2 \emptyset$



This module is an option and must be ordered separately. It allows easy setting of four time ranges in the RUN-mode independently of the program (e.g. for setting delay times).

The repetition accuracy is:

- under constant conditions

Ø. $1 \%$ of the time range set

- under extreme conditions
$1 \%$ of the time range set ( $\mathrm{T}=\emptyset \ldots 5 \emptyset^{\circ} \mathrm{C}, \mathrm{U}=24 \mathrm{~V} \pm 2 \emptyset \%$ )

For waiting for a time to elapse, set at timer 64, use the following simple program:

```
REO }6
SEO }6
```

WIH 64


The corresponding LED lights up while the timer is running down.
Please do not forget that this module is only necessary, if the 32 software timers included in every standard PCAØ do not suffice. The software timers can be modified in the range $\varnothing . \varnothing 1 \mathrm{~s}$ to $65 \emptyset \emptyset \mathrm{~s}$ by the program or via 8 inputs by means of the BCD-Switches in the RUN-mode (see example A8).

### 4.2 The user memory

Three different types of user memories with 4 K (4ø96) user steps each are available:

- RAM-chip 8464 on socket, order no. $4^{\wedge} 5 \emptyset 2^{\prime} 471^{\circ} \emptyset$

This type of memory allows one to write, erase or overwrite a program as desired with the hand-held programming unit $P \emptyset 5$. In case of a voltage failure, the memory contents are stored in the CPU for approximately 2 months thanks to the buffer battery. The program, however, cannot be transported, as it is lost when the RAM-chip is removed.

- Buffered RAM-memory module type PCA1.R95

Contrary to the RAM-chip, the program in this memory can be transported, as it is protected by an integrated electronic system and stored by a lithium battery for approximately 8 years.

Programming with the hand-held programming unit Pø5.

- EPROM-chip 2764 on socket, order no. $4^{-5} 52^{-} 4719^{-} \emptyset$

In an EPROM a program is reliably stored for more than a decade. However, the program cannot be entered directly into the EPROM with the PCAD. For this, the following possibilities are offered (ask for the special documentation):
a) with the EPROM-load module PCA2.P16
b) with the universal programming unit PCA2.P21
c) with the CPUs of the series PCA2 (M31 and M32)

Every EPROM can be erased with an appropriate UV-light source almost as often as desired.

Depending on the user memory in use, the selection jumpers on the CPU must be inserted (see also figure in chapter 5).

*) Position for write-protection
Please note: The jumpers should be repositioned only with the PC switched off.

### 4.3 The programming units

- The hand-held programming unit PCA2. PØ5


This handy programming unit was developed in particular for the series PCAØ, but it can also be used for the series PCA1 and PCA2.

All operating modes can be selected with keys. Programming is performed in the PROG-operating mode by means of a 10 -part keyboard in the easily understandable numerical code. All elements (inputs, outputs, timers, counters) can be interrogated or set in the "MAN"-operating mode.

All timer and counter values can be indicated in the RUN-mode. In the operating mode "STEP" a jump can be effected to any user step of the 4K-user memory. Finally, "BREAK" permits the program execution up to a set break-point and continuation in step-by-step operation. For details refer to "Operating modes" in chapter 8.

- The programming interface РСАП. Рø1


This interface also allows connection of all SAIA-PC programming units with the series PCAØ, namely the following:

- P10 hand-held programming unit with numerical code
- P18 hand-held computer with numerous possibilities
- P21 universal programming unit
- IBM-PC with SAIA-macro-assembler.

As a result, all upwards-compatible members of the SAIA-PC system family are available also for the PCAØ.
5. Brief instructions for operating the РСАФ

a) Function jumpers
(1) When delivered, the function jumpers are inserted as follows:

- Time base " $1 / 1 \varnothing$ " is inserted (for a time base of $1 / 10 \mathrm{~s}$ )
- Flags and registers are non-volatile, when "NVOL" (nonvolatile) is not inserted
- Jumpers for user memory as evident from the above figure apply to the buffered RAM-module PCA1.R95.

If the jumpers are not in these positions, they can be changed with a small screw-driver. In order to provide access to the CPU the cover needs to be removed by two screws.

The jumpers should be repositioned only with the PC switched off.
b) Power supply
(2) Take a transformer (for "playing" $2 \emptyset \mathrm{VA}$ is enough) with a secondary voltage of 24VAC and connect the terminals + and M of the PCAØ via a bridge rectifier. (The PCA2. 5 Ø5 simulation unit already contains this power supply, see section e).
(3) A switch gives the advantage that by switching off the PCAØ all resettable elements and the STEP counter can be easily reset to their initial defined positions.
c) Installation of the user memory R95 and programming unit P95
(4) The buffered RAM-module PCA1.R95 needs to be plugged onto the empty user socket in the specified position (notch on the left).
(5) The programming unit PCA2.Pø5 is connected via the 25 -pole PGUconnector.
If any other programming unit than Pø5 is used, the interface PCAØ.Pø1 needs to be interposed.
d) Program example "Blinker"
(6) Switch on voltage supply. Yellow lamp "CPU RUN" blinks every $2 s$ (1s on, is off).
(7) Select the operating mode "PROG" by pressing the $P$-key of the programming unit (for at least $\emptyset .5 \mathrm{~s}$ ). As a result, the red LED "PROG" on the P 95 lights up.
(8) Enter the following blinker program:

|  | STEP | CODE | OPERAND | Program in mnemonic code |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $A, \emptyset, E$ | ( $\emptyset \emptyset \emptyset)$ | (ø) | ( $00 \emptyset 0)$ |  |  |
| E | (øøロ1) | $\emptyset 2$ | 256 | $\rightarrow$ STL | 256 |
| E | ( 002 ) | 14 | 256 | STR | 256 |
| E | (øøø3) | $\emptyset \emptyset$ | 5 |  | $0.5 s$ |
| E | (øØ04) | 13 | 24/4Ø ** | COO | 24/4ø ** |
| E | (øø05) | 20 | 1 | - JMP | 1 |
| E | (øøø6) | ( $\varnothing \square)$ | ( $0 \square \square \square)$ |  |  |

[^1](9) Set program counter to zero:

Key sequence $S$ operating mode STEP, as confirmation the red LED "STEP" lights up
(A) address, $\square$
(10) Select operating mode "RUN":

Press key [R (RUN) for $\emptyset .5 \mathrm{~s}$
$\longrightarrow$ Red LED "RUN" lights up on Pø5
$\longrightarrow$ Green lamp "RUN" lights up on PCAø
$\longrightarrow$ Output 24 or $4 \emptyset$ blinks $\emptyset .5 s$ on and $\emptyset .5 \mathrm{~s}$ off (frequency 1 Hz )
e) Connection of the input simulation unit PCA2.S $\$ 5$

Including the I-simulation unit PCA2.Sø5 gives a complete set of programming and practicing devices which can be used to try out all examples of programs contained in this manual.


Instead of the I-simulation unit PCA2.S05 the bigger PCA2.S10 can also be used with the connecting cable PCA1.K8Ø.

6. Detailed information on power supply and watchdog
6.1 Power supply of the PCAØ

Supply voltage $U_{\text {in }} \quad$ 24VDC smoothed or pulsating
Tolerance for $U_{i n}$

- general
- for version with relays
$\pm 2 \emptyset \%$
pulsating voltage $\pm 2 \emptyset \%$
( $\mathrm{T}_{\mathrm{amb}}=\emptyset \ldots 5 \emptyset^{\circ} \mathrm{C}$ )
smoothed voltage $+2 \varnothing \%$
( $\mathrm{T}_{\text {amb }}=\emptyset \ldots 35^{\circ} \mathrm{C}$ ) $-5 \%$
Supply current
- PCAD.M12T
- PCAD.M14R

Fuse
max. 0.5 A (with $\mathrm{P} \emptyset 5$ connected)
max. 0.9 A (with $\mathrm{P} \varnothing 5$ connected)
1.6A quick-acting

Several components protect the PCAØ against interference voltages, wrong polarity and voltage drops. The 5 V for supplying the electronic components is generated by means of a switching regulator.


### 6.3 The WATCHDOG-monitoring circuit

The WD-circuit reliably monitors the correct execution of the user program. In case of an error effective safety measures can be taken.

The WD-relay remains excited (contact H-J is closed) as long as the address 255 receives an alternating signal of $\geq 5 \mathrm{~Hz}$. This signal is generated in a circulating program simply with the instruction COO 255. During normal operation of the CPU in the RUN-mode, the terminals $\mathrm{H}-\mathrm{J}$ remain closed and the green WD-lamp lights up. If a malfunction occurs in the CPU or any other operating mode than "RUN" is selected, the contact $\mathrm{H}-\mathrm{J}$ opens, the WD-lamp goes out.

For critical systems it is recommended to make use of the WD-function in combination with the following safety circuit. Upon releasing of the WDrelay the PCAD is no longer supplied with voltage, which has the result that all outputs are reset at once.

Contact rating of the WD-contact: $1.5 \mathrm{~A}, 48 \mathrm{VAC}$ or DC

7. Detailed information on inputs and outputs
7.1 Inputs 24VDC (B9ø)

Input voltage $U_{i n}$
Voltage level

Input current at 24VDC
Input delay
Operating mode

Source operation

24VDC, smoothed or pulsating
H: 19...32VDC
L: Ø... 4VDC
$1 \not ⿴ \mathrm{~mA}$
9 ms
Source or sink operation, depending on the connection

(Terminal assignment)

*) The inputs 16... 19 of the PCAØ.M12T can be acted upon only in source operation.
(Terminal assignment)


### 7.2 Transistor outputs (B9ø) Type T

Output current range

Voltage range $U_{\text {out }}$
Voltage drop
Operating mode

5mA. . . 0.5 A
In the voltage range 5...24VDC the load resistor should have a value of at least 480 hm
5...36VDC, smoothed or pulsating

Max. 1.5 V with $\mathrm{I}=\emptyset .5 \mathrm{~A}$
Source operation (positive switching)

(Terminal assignment)

7.3 Relay outputs (A21) Type R

Type of contact
Contact rating

Contact protection
Contact life (AC 1)

Pure silver, make contact
$3 A, 25 \emptyset V A C$ AC1
1A, 250 VAC AC11
3A, 24VDC DC1*
1A, 24VDC DC11*
3.3 nF and 330 hm

3A, 22øVAC: number of switching cycles is $\emptyset .1$ mio.
1.5A, 22ØVAC: number of switching cycles is $\emptyset .5$ mio.
$\emptyset .3 \mathrm{~A}, 22 \emptyset \mathrm{VAC}:$ number of switching cycles is 5 mio.
(Terminal assignment)


*) For reasons of life and reliability transistor outputs should be preferred for switching a DC-voltage of 24 V .
8. The operating modes of the PCAD

In order to prepare and test programs, it is necessary to operate the PC in different operating modes. This is effected by pressing the corresponding keys on the small programming unit PCA2.Pø5 or with the programming interface PCAØ.PØ1 for other programming units.

Please note that these operating mode keys must be actuated for at least 0.5 s for reasons of security. The actual selection of the appropriate operating mode is confirmed by the indicating LED of the programming unit Pø5. When the programming unit connector of the PCAD is disconnected, the selected operating mode is maintained.

When switching on the PCAØ, the following operating modes are selected automatically:

- With the programming unit connected $\rightarrow$ STEP (the LED "STEP" on the Pø5 lights up, the green LED "RUN" on the PCAØ does not light up!).
- Without programming unit $\rightarrow$ RUN (the green LED on the PCAØ lights up).
8.1 Operating modes (summary)

R RUN Normal program execution (lamp RUN on the PCAØ lights up)
P PROG A user program can be loaded into a RAM-memory (plugged onto the user socket of the PCAØ).

M MAN Manual interrogation and setting of elements (inputs, outputs, flags, timers, counters).

S STEP Jump to a preselected step address of the user program and step-by-step execution.

B BREAK Program execution up to a set "breakpoint" and subsequent step-by-step operation.
$T$ TEXT Has no function in the standard PCAD.

### 8.2 Detailed description of the operating modes

(A) RUN Normal program execution

The PCAD is automatically in the RUN-mode when switching on if the programming unit is not connected.
$P$
PROG Programming
A program can be stored in a RAM-memory (on the user socket of the PCAØ) or overwritten (corrected).

( + Terminates the input
Test program $\quad+\square$ or $-\square$
M MAN ** Manual interrogation or setting of elements
(Elements = inputs, outputs, flags, counters, timers)

Interrogation:
 display of the logic state

Element address
Setting: $\quad A] \underbrace{x \times x}[E] \rightarrow$ or $\square$
Element address
$\square$ STEP $+\rightarrow$ Display showing where the program is.
Jump to the preselected step address of the user program
[A] $139 \square$ Program jumps to step 139, then
$\pm \pm \ldots$ step-by-step execution of the program with the
result of the logic operation being checkable 㭗 $A C C=1$ :
Switching to RUN is always possible.
In case of parallel programs, only the activated parallel
program is executed in the STEP-mode
B BREAK Interruption of the program run and subsequent step-by-step operation
$\square \longrightarrow$ Display showing where the program is
$\square]$
... step-by-step execution of the program with the
result of the logic operation being checkable * $A C C=1$ *
Switching to RUN is always possible.
In case of parallel programs, all programs are executed simultaneously (as in the RUN-mode).

Setting of a breakpoint
(A) $82 \emptyset \square$ program runs up to step 820 , then
$\pm+\ldots$ step-by-step operation skipping the "critical" point.
*) $\mathrm{ACC}=$ accumulator is used to indicate the result of the logic combination. If $A C C=1$, (conditions of the logic combination fulfilled), the following switching instructions are executed.
**) If the address of a timer or counter is preceded by a 3 (e.g. $326 \emptyset$ for counter $26 \varnothing$ ), the value of this register can be read or entered manually with $[\mathrm{E}$ value $\square$.
9. Programming in three easy steps

> Combined programming
> C incl. arithmetic and index register

## B Programming according to flow-chart

A Programming according to a ladder diagram

The PCAD is equipped with a very efficient instruction set, level (1H). Thanks to this instruction set even complex control problems can be solved easily. The programs of the РСАП can be used at any time with other series of the SAIA-PC system family, too. This enables your controller to "grow" according to your requirements, without having to write new programs every time.

In order to make the start of programming easier for the beginner, the performance of the PCAD is split up into three easily understandable steps. Maybe your control problem is so easy that it can be solved even on level A.
Let's start programming with a ladder diagram, although this is no longer suitable for modern programming, because process runs are often particularly difficult to represent in relay logic.
But you will see that your PCAØ understands any ladder diagram.
(A) Programming according to ladder diagram


The program prepared according to a ladder diagram is entered into the user memory (RAM) by means of the programming unit. In the RUN-mode, the CPU reads this program line by line and checks the relevant inputs. If one of the contacts I2 or I4 is closed, an "H" (High = voltage greater than +19 V ) is stored in the CPU and after reading the 2nd line it is ANDconnected with the latter. If both contacts are closed (H), the state of the accumulator $=1(\operatorname{ACCU}=1)$ and the output 27 is activated, the contactor at output 27 is actuated.

In line 4 the CPU processes input 7. If this contact is closed, output 24 is activated, the indicating lamp lights up.

We can combine more of these logic operations, because our user memory of $4 \mathrm{~K}=4096$ program lines is extremely large. If all logic operations have been programmed, we have to tell the CPU to return to line 1. In this way, our program is permanently run cyclically at a high speed, and all alterations are immediately transferred as logic operation results from the inputs to the outputs.

## Instruction set (A) ladder diagram

The instructions available for programming such logic operations are classified into logic instructions and switching instructions.

Their functions are listed in the following table. Do not let yourself be confused by the large number of functions. When programming according to a ladder diagram only about 8 of them are used frequently, which you will soon know by heart.


Each instruction in the user program consists of 1 line (in certain cases of 2 lines). In addition to the line number or step address (STEP) a line contains the instruction code (CODE) and operand (OPERAND). The instruction code indicates "WHICH" instruction is to be executed, and the operand determines "WHERE" this instruction is executed.

Structure of the program or instruction line:


STEP The line no. defines the position of the instruction in the user memory. Decimal numbering from Ø...4095 (4K).

CODE Depending on the programming unit the instruction code can be entered in a 3 -digit mnemonic code or in numerical code from $\emptyset$ to 31. The mnemonic code is based on abbreviations of the corresponding English instructions. It is therefore easy to remember and understood internationally.

OPERAND Here, the address of an element (input, output, timer, counter or flag) or in case of jump instructions the destination address (line no.) is entered.

Timing and counting instructions consist of two lines. In the second line, the appropriate time or counter value appears in the operand.
(A1) A first programming example
Before starting to enter the program, it must be noted down in mnemonic code on paper. For this, it is advantageous to copy the programming lists added at the end of this manual.


For programming and simulation we establish the same configuration as described in section 5. All the following examples always refer to the small PCAQ with only 16 or $2 \emptyset$ inputs. For the bigger versions the output addresses are added in brackets respectively.

With the programming unit PCA2.Pø5 the above program can now be entered into the plug-in user memory (RAM or R95).
$P \longrightarrow$ operating mode PROG
(A)
2)
3)
4)


STEP Mnemonic Numerical OPERAND code code

| $\emptyset \emptyset \emptyset$ | (NOP) | $\emptyset \emptyset$ | $\emptyset \emptyset \emptyset \emptyset$ |
| :---: | :---: | :---: | :---: |
| 1 | (STH) | $\emptyset 1$ | 2 |
| 2 | (ANH) | $\emptyset 3$ | 4 |
| 3 | (OUT) | $1 \emptyset$ | 27 |
| 4 | (STH) | $\emptyset 1$ | 8 |
| 4 | (STH) | $\emptyset 1$ | 7 |
| 5 | (OUT) | $1 \emptyset$ | 24 |
| 6 | (JMP) | $2 \emptyset$ | 1 |

1) Selection of the step address $\emptyset$ ( $A]=$ Address) and erasure of the contents with [E] = Enter.
2) Every following E increments the step address by 1 and the contents of the old program line are erased and prepared for the new input.
3) 8 was entered accidently instead of 7 . Correction with $C$ and repetition of the instruction. The step address is not incremented as a result of [C].
4) The last input must be stored with $\square,[E, A]$ or $\square$.

After entering the program it is recommended to compare the program stored in the user memory step by step with the program previously written down:

## (A) $1+\square+\ldots$

Now of course we want to test whether the program runs.

S $\longrightarrow$ The LED indicating operating mode "STEP" lights up
A $1 \square \longrightarrow$ The program execution should start at step 1
R $\longrightarrow$ The LED "RUN" lights up, the program is running

Now close contacts I2 and I4 $\longrightarrow$ the LED of 027 lights up. Upon closing I7 $\longrightarrow 024$ lights up.

Ladder diagram


Program

| Step | $\begin{aligned} & \text { Mnemo. } \\ & \text { code } \end{aligned}$ | $\begin{array}{\|l\|l\|} \substack{\text { Numer ic. } \\ \text { coder }} \end{array}$ | Operand | Comment |
| :---: | :---: | :---: | :---: | :---: |
| F-11 | STH | 01 | 3 | Interrogation 13 |
| 11 | ANH | 03 | 2 | AND 12 |
| 12 | ORH | 05 | 6 | OR I6 |
| 13 | ORH | 05 | 4 | OR I4 |
| 14 | ANH | 03 | 1 | AND If |
| 15 | ANH | 03 | 7 | AND I7 |
| 16 | OUT | 10 | 27 | Output 027 |
| 17 | OUT | 10 | 26 | Output 026 |
| -18 | JMP | 20 | 10 | Retuan |

## Note:

- Every OR-instruction starts a new branch of the parallel connection at the very beginning on the left. Afterwards, AND-operations can be added again to this parallel branch. The whole program, however, is considered as only one logic operation.
- If the logic operations are finished, as many actions as desired (depending on this logic operation) can be added.
- Enter the program as follows:
$P \longrightarrow$ LED "PROG" for programming mode lights up

A 1 1 $\square$ 国 $\longrightarrow$ The PCAØ is prepared to accept the above program from step address 10

Continue to proceed as described in example A1, but make the PCAD execute the program from step 10 on with $S$, $A$, $+\pi$. Run!
(A3) OR is "stronger" than AND and getting to know other "elements"

The following contact arrangement must be programmed:


One is tempted to prepare the program as above. However, as $O R$ is always at the beginning of a parallel branch, the following circuit would be the result:


There are two possibilities of implementing the desired function:
a)

$\left[\begin{array}{rllr}2 \emptyset & \text { STH } & \emptyset 1 & 6 \\ 21 & \text { ANH } & \emptyset 3 & 2 \\ 22 & \text { ORH } & \emptyset 5 & 6 \\ 23 & \text { ANH } & \emptyset 3 & 3 \\ 24 & \text { OUT } & 1 \emptyset & 27 \\ 25 & \text { JMP } & 2 \emptyset & 2 \emptyset\end{array}\right.$

The contact I6 is programmed in both parallel branches.
b)

$\left[\begin{array}{rrr}2 \emptyset & \text { STH } & 2 \\ 21 & \text { ORH } & 3 \\ 22 & \text { OUT } & 24 \\ 23 & \text { STH } & 6 \\ 24 & \text { ANH } & 24 \\ 25 & \text { OUT } & 27 \\ 26 & \text { JMP } & 2 \emptyset\end{array}\right.$

Programming is performed in two steps. As evident from possibility b) outputs can be used as desired in other logic operations, too.

It would be a pity to sacrifice an extra output for this easy task. Therefore, every PCAØ contains 712 FLAGS!

The elements and registers of the PCAD:

Elements
ADDR.

$\vdots$
288
32 counters C*, which can also be used as
flags

$\left[\begin{array}{cc}\vdots & \vdots \\\right.$\cline { 1 - 1 } \& 256\end{array}$\} 32$ timers $T^{*}$ or counters $C^{*}$
$\left[\begin{array}{cc}\vdots & \vdots \\\right.$\cline { 1 - 1 } \& 256\end{array}$\} 32$ timers $T^{*}$ or counters $C^{*}$
$\left[\begin{array}{cc}\vdots & \vdots \\\right.$\cline { 1 - 1 } \& 256\end{array}$\} 32$ timers $T^{*}$ or counters $C^{*}$


## 67 66 4 hardware timers, provided

 65 that the module PCAØ. $\mathrm{H} 2 \varnothing$ 64 is attached
$\left.\begin{array}{l}63 \\ :\end{array}\right\}$ max. 64 inputs I and outputs 0

Register
ADDR.

| 319 | 64 registers* <br> as counters or <br> timers at 16 bit |
| :--- | :--- |
| 256 |  |

*) By inserting the jumper "NVOL" on the CPU all of these locations can be made non-volatile, i.e. when switching off the PCA $\emptyset$ these data are not lost.

With the aid of a flag the ladder diagram can now be drawn as follows:


| ADDR | NC | MNC | OPRD |  |
| :---: | :---: | :---: | :---: | :---: |
| -3ø | $\emptyset 1$ | STH | 2 | Interrogating 12 |
| 31 | $\emptyset 5$ | ORH | 3 | OR I3 |
| 32 | 10 | OUT | $40 \emptyset$ | Storing the intermediate result on flag $4 \emptyset \emptyset$ |
| 33 | 01 | STH | 400 | Interrogating flag $4 \emptyset \emptyset$ (and thus the OR-function) |
| 34 | $\emptyset 3$ | ANH | 6 | AND 16 |
| 35 | 10 | OUT | 27 | Output of the result via 027 |
| -36 | $2 \emptyset$ | JMP | $3 \emptyset$ | Return to the beginning |

Two kinds of start/stop circuit with latching contactor
a) Presented in a ladder diagram

The following classical example is known from the technique of contactors:


Program
$\left.\begin{array}{cccc}\text { ADDR } & \text { NC } & \text { MNC } & \text { OPRD } \\ 4 \emptyset & \emptyset 1 & \text { STH } & 1 \\ 41 & \emptyset 5 & \text { ORH } & 24 \\ 42 & 1 \emptyset & \text { OUT } & 4 \emptyset 1\end{array}\right\}$ Start

As in example A3 programming is performed using a flag. As evident from the program, the normally open contact I $\emptyset$ is connected with " H ", as 024 can be activated only with contact iø being closed.

This way of programming is fail-safe against wire break. If a wire breaks in the lines of $1 \emptyset$, I1 or 024,024 is always inhibited.
b) Logic diagram presentation

## Program


$\left.\begin{array}{|cccc}\text { ADDR } & \text { NC } & \text { MNC } & \text { OPRD } \\ {\left[\begin{array}{cccc}5 \emptyset & \emptyset 1 & \text { STH } & 1 \\ 51 & \emptyset 3 & \text { ANH } & \emptyset \\ 52 & 11 & \text { SEO } & 24\end{array}\right\} \text { Start }} \\ \hdashline 53 & \emptyset 2 & \text { STL } & \emptyset \\ 54 & 12 & \text { REO } & 24 \\ 55 & 2 \emptyset & \text { JMP } & 5 \emptyset\end{array}\right\}$ Stop

SEO (11): Set Output With this instruction an element (output or flag) is continuously set until it is reset with REO.

REO (12): Reset Out- With SEO/REO we are able to program a flip-flop. put Both instructions are executed only if the result of the logic operation is positive ( $A C C U=1$ ).

The "Set" instruction of example b) is executed only if I $\emptyset=H$. If both keys are pressed, the "Reset" instruction takes precedence because of the AND-operation.

This way of programming is also fail-safe against wire break.

Example of a pulse divider (stepping switch) using the instruction DYN and COO

The following function is referred to as a pulse divider (stepping switch):


The output 027 is modified only upon each rising edge of I3.

Presenting this function in a ladder diagram results in a quite complicated diagram and a long program.
a) Ladder diagram presentation

Try once on your own to draw a relay diagram that realizes this function. It will not be easy!
b) With SEO/REO and the above function drawn as a logic diagram, it is already much easier.

c) With the instructions DYN and COO

DYN (99): Dynamic execution of an interrogation function or signal edge triggering. With the DYN-instruction (together with a flag in the operand) the preceding interrogation instruction accepts only the positive alteration (rising edge). A permanent H -state or the falling edge are ignored.

COO (13): Complement output. The logic state of an output or flag is tested and inverted i.e. if an output is set, it is reset with COO and vice versa.

With these efficient instructions this problem can be solved really easily.


| ADDR | NC | MNC | OPRD |
| ---: | ---: | ---: | ---: |
| $\left[\begin{array}{r}6 \emptyset \\ 61\end{array}\right.$ | STH | 3 |  |
| 61 | $\emptyset 9$ | DYN | $5 \emptyset \emptyset$ |
| 62 | 13 | C00 | 27 |
| 63 | $2 \emptyset$ | JMP | $6 \emptyset$ |

Interrogation of I3
Edge triggering, storing in flag $50 \emptyset$ Testing the state of 027 and negating it

Without the DYN-instruction in this example the output 027 would be complemented in each program cycle, if switch I3 were closed; that is, approximately $3 \varnothing \emptyset$ times per second in this short program loop.

With DYN, output 27 will be complemented only in the 1st program cycle upon closing of I3. Every other cycle will not have any effect, until the signal state of I3 has changed and a new rising edge is formed.
(A6) The software timers with switch-off delay
In example A3 all elements of every standard PCAØ and their corresponding addresses were shown. The 32 software timers reside on the addresses 256...287. For each address there is a 16 -bit register in which values from Ø... 65535 can be stored.

For setting and starting a software timer a two-line instruction is necessary:

| $\operatorname{STR}(14)$ | $256:$ |
| :---: | :---: |
| $\emptyset \emptyset$ | $2 \emptyset:$ |

Starting the timer with address 256
Loading the corresponding register with the value $2 \emptyset$. (Direct input up to $2 \emptyset 47$ for code $\varnothing \varnothing$ ).


When the CPU reads this instruction, the logic state of timer 256 is set to " H " and the register value is reduced according to the time base $(1 / 1 \emptyset s)$. After this has been done $2 \emptyset$ times ( $2 \emptyset \times 1 / 1 \emptyset s=2 s$ ) the logic state of the timer 256 is "L" again.

With this basic function a switch-off delay can be implemented easily:


| ADDR | NC | MNC | OPRD |
| :---: | :---: | :---: | :---: |
| $\rightarrow 70$ | $\emptyset 1$ | STH | 7 |
| 71 | 14 | STR | 256 |
| 72 | $\not \varnothing$ | $\emptyset \emptyset$ | 75 |
| 73 | $\emptyset 1$ | STH | 256 |
| 74 | $1 \varnothing$ | OUT | 27 |
| -75 | $2 \emptyset$ | JMP | $7 \emptyset$ |

Interrogation of I7
Set timer and start
Input of time in $1 / 10 \mathrm{~s}\}$
2-l ine instruction
Interrogation of timer T256
Transfer to 027

Upon closing of 17 the timer is set and its logic state is set to " H ". The time does not start to run down before I7 is opened. (In fact, the time starts to run down at once. But as the timer is set again in the next program cycle after some $1 \emptyset \emptyset \mu \mathrm{~s}$ with I 7 closed, the time starts to run down again from the very beginning, until the signal present at I7 is removed, i.e. the contact 17 opens).

If I7 is closed again while the time is running down, the timer is reset and started again.
P.S.: As evident from chapter 5 the time base can be reduced to $1 / 1 \emptyset \emptyset \mathrm{~s}$ by removing the jumper " $1 / 1 \varnothing$ " on the CPU thus enhancing the resolution.

## รaiヨ’ロㄷ

(A7) Use of the hardware timer module PCAØ. H2 $\emptyset$
In addition to the 32 precise software timers (which are included in every standard version) 4 hardware timers (additional module) are available in order to use time functions in the programs.

The 4 hardware timers with the addresses $64 \ldots 67$ can be operated in a similar way to the software timers.

Sof tware timers

| STR | 256 |
| ---: | ---: |
| $\emptyset \emptyset$ | t |

Hardware timers

| REO | 64 |
| :--- | :--- |
| SEO | 64 |

For the hardware timer the time range and time are set directly on the hardware module PCAØ. H2Ø.

Example according to A6:


As with the software timer, the time starts to run down at once in case of the HT 64, too. As long as 17 is switched on however, HT 64 is reset in each program cycle. The time actually starts to run down when 17 is opened resulting in the above switch-off delay.
(A8) Software timers with fleeting-on delay with external time input in BCD-code

The PCAØ also allows reading the BCD-values directly into the registers of the timers and counters. As a result, time values can be changed at any time with the BCD-switch. If we use the input simulation unit PCA2.SØ5 as described in chapter 5, we can see that the input addresses I8...I15 are acted upon by a two-digit BCD-switch. For this, the BCDswitch transmits the following signals to the input:


Now, the STR-instruction can expanded in such a way that the delay time can be directly read off the BCD-switch.

Code OPERAND BCD-switch

| STR (14) | 256 |
| :---: | ---: |
| NC | En) |

1st line 2nd line

(En) : The highest input address of the two BCD-switches (e.g. (15)) is introduced into the operand.

NC : The numerical code has the values 16,17 or 18 with the following meanings:

$$
\begin{aligned}
16 & \text { value } \emptyset \ldots 99 \times \text { time base }=\emptyset \ldots 9.9 s^{*} \\
\text { or } 17 & \text { value } \emptyset \ldots 99 \emptyset \times \text { time base }=\emptyset \ldots 99 s^{\star} \\
\text { or } 18 & \text { value } \emptyset \ldots 99 \varnothing \emptyset \times \text { time base }=\emptyset \ldots 99 \emptyset s^{*}
\end{aligned}
$$

[^2]Problem:
The time range 1 to 99s must be set with the external BCD-switch. The inputs 18... 15 are acted upon by the two BCD-switches. The time function is to have a fleeting-on delay.

Because of the DYN-instruction only the rising edge of $I \emptyset$ is taken into account, enabling the timer 258 to run down without interruption.


Program:


If the same function has to be performed by a hardware timer, just replace the lines 92 and 93 as follows:

(A9) Switch-on delay with two new, very useful instructions
Problem:


In case of the switch-on delay the timer is also started by the rising edge of I7 as in case of the fleeting-on delay. However, output 27 should be switched on only after the time has elapsed. order to activate 027, the rising edge of 17 must have been generated and the timer run down (L). This can be shown by means of the following logic diagram:


As we are working with an intelligent PC, elapsing of the time has to be displayed on the programming unit. This is achieved with the instruction DTC (31): "Display Timer or Counter". Provided that this instruction is executed at least once every second (which does not constitute a problem in circulating programs), the actual contents of the corresponding timer or counter are displayed in the operand field of the P $\varnothing 5$. Activation of the DTC is effective only if the $A C C U=1$. Therefore, it needs to be pre-


Program:

| ADDR | NC | MNC | OPRD |  |
| :---: | :---: | :---: | :---: | :---: |
| -100 | $\emptyset 1$ | STH | 7 | Interrogation of 17 |
| 191 | 09 | DYN | 402 | Signal edge triggering |
| 102 | 14 | STR | 287 | Start timer and set it to |
| 193 | $\emptyset \emptyset$ | $\emptyset 0$ | 120 | $12 \emptyset \times \emptyset .15=12 s$ |
| 104 | $\emptyset 1$ | STH | 402 | Interrogation of the edge flag of I7 |
| 105 | 94 | ANL | 287 | AND timer run down (L) |
| 196 | 10 | OUT | 27 | then output to 027 |
| 107 | 19 | SEA | $\emptyset$ | Set $\mathrm{ACCU}=1$ |
| 198 | 31 | DTC | 287 | Display of the timer contents |
| -109 | 20 | JMP | 100 |  |

If this function is to be performed by the hardware timer HT 67, the lines 1ø2/1ø3 and $1 \varnothing 6$ must be replaced as follows:

| ADDR | NC | MNC | OPRD | Start hardware timer HT 67 |
| :---: | :---: | :---: | :---: | :---: |
| 192 | 12 | REO | 67 |  |
| 193 | 11 | SEO | 67 |  |
| - |  |  | - |  |
| 106 | 04 | ANL | 67 | AND HT 67 run down (L) |

(B)

Programming according to a flow-chart
Show a ladder diagram and this flow-chart to a 12-year old pupil. What do you think will he interpret?


So, is it surprising that more and more industrial processes, too, are being described using a flow-chart. Especially in the chemical or foodprocessing industry, as well as in many branches of mechanical engineering, there are many processes that can be described in an easy and understandable way with the aid of a flow-chart.

The PCAØ comprises many instructions which can be used efficiently for programming according to a flow-chart resulting in much easier and clearer functions.

Counter instructions can be used of course both for programming with the aid of a ladder diagram and flow-chart. In order to prevent level (A) from getting too complex, we have assigned them to level (B).

(B1) Upwards/downwards movement

## Problem:

As a result of a pulse at $I \emptyset$ the load L must be moved upwards ( $U P=027$ ) until I4 is opened. Then, the load must be lowered again (DOWN $=024$ ) until I7 is opened.

Additional conditions:

- If Iø is permanently on this sequence must be continuously repeated.
- After a voltage break an upwards movement must not be triggered before a pulse is present at $1 \varnothing$.
- Note: Before starting the program the switches I4 and I7 must be closed.


Problem solving by means of a flow-chart:

Flow-chart


Comment

Wait as long as $I \emptyset$ is open (low).
If $I \emptyset$ is closed, then ...

Set output 27 (UP)

Wait as long as I4 is closed (high).
If I4 is opened, then ...

Reset 027 (STOP) and set 024 (DOWN)

Wait as long as 17 is closed (high).
If I7 is opened, then...

Reset 024 (STOP)

Return to the beginning

As evident from the example, a sequence is described step-by-step when programming according to a flow-chart. The stages of the process are split up into conditions (wait for an input state) and actions (set or reset outputs).

If we follow this program in step-by-step operation in the operating mode STEP (key sequence ( $S$ ( $A 110+\square$ )...., we will notice that the processor itself stays in the wait loop, until the condition for continuing is fulfilled. This means: the program does not permanently circulate cyclically like the program prepared with the aid of a ladder diagram, but the program is executed according to the progress of the process.

This has three important advantages:

- The programs are made clearer, because they show the individual steps of a process and not an abstraction in a ladder diagram which is not related to the process.
- Only those program sections are processed which are of importance for the current stage of the process. As a result, possible malfunctions resulting from the execution of irrelevant program parts are avoided. Most importantly, however, the reaction time between step condition and action is considerably reduced.
- If the slide L stops after the first up-/down-movement, it is possible to exactly localize the error in the program by means of the programming unit in the operating mode "STEP", $+{ }_{+}+\ldots$ At step $11 \emptyset$, WIL $\emptyset$ is displayed. Moreover, the LED of I $\emptyset$ indicates that this input has not been activated. Therefore, the error can be quickly narrowed down to the contact or the connecting line of Iø.


## (B2) Up-/downwards movement with timer and counter

## Problem:

The movement described in example B1 is to be effected not only once but a certain number of times (e.g. 5 times) and a pause of 2.5 s must be made each time at the points where the direction is reversed.

Solution:


Contrary to problem B1, in this example timers were also integrated into the action parts. One must wait for these to run down in additional conditional loops (WIH 256).

The use of the instruction SCR (15): Set counter is new.
This is a two-l ine instruction (as in case of the timer). In our example the counter is set to 5 . As long as the counter C3ø is greater than zero, its logic state is "H".


With the instruction DEC (18): Decrement Counter the counter state is decremented by 1 in each run.

The interrogation instruction STH $30 \emptyset$ allows to continuously check the logic state of the counter. If the $A C C U=1$ after the interrogation, the counter contents are still greater than zero and another loop must be executed. Returning to the beginning of a loop is effected with the conditional jump instruction JIO (21): Jump if ACCU $=1$.

In other words: If the counter is greater than zero, another loop must be executed: If the counter state is zero, no more jumps are effected and the last JMP-instruction leads back to the program beginning.

If the hardware timer HT 65 is used instead of software timer T256, the corresponding program parts need to be replaced as follows:
$\left.\begin{array}{rr}\text { STR } & 256 \\ \emptyset \emptyset & 25 \\ \text { WIH } & 256\end{array}\right\} \longrightarrow\left\{\begin{array}{rr}\text { REO } & 65 \\ \text { SEO } & 65 \\ \text { WIH } & 65\end{array}\right.$
(B3) What needs to be done, if we want other functions to be performed together with the upwards/downwards movement?

In this case, we prepare a second or third parallel program out of the 16 parallel programs provided for each PCAD-standard version!

Problem:
Simultaneously and independently of program B2 a blinker with a flashing time of 0.3 s has to be generated in a parallel program and the actual counter state of the upwards/downwards movement has to be continuously displayed on the programming unit.

Program:


In the very beginning, we use the two-l ine instruction PAS (29) 1 (program assignment) to assign the number (1) to the parallel program starting at step address 130 in the so-called assignment part.

Other parallel programs can be added easily with PAS (2) or PAS (3).
The program at step addresses 123 to 129 is a so-called circulating program without wait loops. It may include other functions such as monitoring or other permanently performed tasks. This parallel program without assignment is automatically given the number ( ( ).

If you want to execute PPØ at the same time from step address 123 on and PP1 from step address 130 on, then jump to the start address $12 \emptyset$ with the following assignment: $S$ (A) $12 \emptyset \quad \square$.

Now, test both programs! Both are executed separately and asynchronously. As mentioned above, up to 16 parallel programs of any desired length can be executed with the PCAØ (as with all SAIA-PCs).
(B4) Subroutines help you save much time and result in shorter and clearer programs
Problem:
The following sequence has to be executed after closing of I1:


Solution a: Of course with the aid of a flow-chart (if you want, try to prepare the program with the aid of a ladder diagram!).


The flow-chart shows the program sequence without subroutines. The program parts which are marked with the brackets are repeated 6 times. Therefore, it is of advantage to write them as subroutines.

Solution b: with subroutine
Main program
(A)

(B)

(c)

(D)

(F)


Thanks to the instruction JMS (23): Jump to Subroutine the same program parts must be prepared only once. Every subroutine ends with the instruction RET (24): Return with the operand $\emptyset$.

| $=================$ Subroutine 182 (wait 1s) |  |  |  |
| ---: | ---: | ---: | ---: |
| $\Rightarrow 182$ | 14 | STR | 256 |
| 183 | $\emptyset \emptyset$ | $\emptyset \emptyset$ | $1 \emptyset$ |
| 184 | 25 | WIH | 256 |
| 185 | 24 | RET | $\emptyset$ |

It is also possible to program the subroutine of the subroutine of the subroutine i.e. down to the 3 rd level.
(C) Indexing, arithmetic and check sum

At the programming level C you are already a professional!
Certainly, one can live happily without climbing up to this level. But, once you are up here, you will be proud of yourself and the РCAØ for having prepared your programs in such a nice and efficient way.

These are the remaining instructions for the software level (1H) of the SAIA-PC system family.

(C1) In case of short programs address indexing has considerable effects
Problem:
All outputs of the PCAØ are to be activated upon closing of input I $1 \varnothing$.


Without address indexing this program would consist of 34 steps for a PCAØ with 32 outputs.

Solution:
With indexing, however, the program will always consist of only 6 lines, irrespective of the number of outputs.

| ADDR | NC | MNC | OPRD |  |
| ---: | ---: | ---: | ---: | :--- |
| $-2 \emptyset \emptyset$ | 16 | SEI | $\emptyset$ | Set index register (IR) to starting value $\emptyset$ |
| $-2 \emptyset 1$ | $\emptyset 1$ | STH | $\emptyset$ | Interrogation of input $\emptyset$ |
| $2 \emptyset 2$ | $1 \emptyset$ | OUT | $1 \emptyset 16$ | Set indexed outputs |
| $2 \emptyset 3$ | 27 | INI | 15 | Incrementing of IR up to the final value 15 |
| 20 | 21 | JIO | $2 \emptyset 1$ | Return, as long as IR $<15$ |
| $2 \emptyset 5$ | $2 \emptyset$ | JMP | $2 \emptyset \emptyset$ |  |

In the following flow-chart the individual functions are easier to understand:

Indexing loop


With the instruction SEI (16) the index register is set to $\emptyset$ (max. value is 255).
$1 \emptyset \emptyset \emptyset$ is added to the first output address 16. As a result, this indexed instruction is processed by the CPU as follows: 1016 - $190 \emptyset$ + IR

As a result of instruction INI (27) the index register is incremented by 1 in each run until it reaches the value 15 .

As long as $I R<15$, the $A C C U=1$, which causes the return.
If $I R=15$, the indexing loop is left. All outputs from $16 . . .31$ have been processed.
(C2) Thanks to upward and downward indexing high flexibility is obtained
Problem:
Upon closing of ID the outputs 016... 31 (63) are to switch on one after the other every $\emptyset .3 \mathrm{~s}$. Upon opening of $1 \varnothing$ the outputs 031 (63)... 16 should be switched off again in reverse order every $\emptyset .1 \mathrm{~s}$.

Solution:


Wait while
contact $\emptyset$
is closed


| ADDR | NC | MNC | OPRD | TURN-ON PHASE |
| :---: | :---: | :---: | :---: | :---: |
| $21 \emptyset$ | 26 | WIL | $\emptyset$ |  |
| 11 | 16 | SEI | $\emptyset$ |  |
| $-212$ | 11 | SEO | 1016 | $016+100=1016$ |
| 13 | 14 | STR | 260 |  |
| I 14 |  | $\emptyset \emptyset$ | 3 |  |
| 115 | 25 | WIH | $26 \emptyset$ |  |
| 16 | 27 | INI | 15 | 031-16 = 15 |
| - 217 | 21 | JIO | 212 |  |
|  |  |  | -- | TURN-OFF PHASE |
| 218 | 25 | WIH | $\emptyset$ |  |
| 19 | 16 | SEI | 15 | 031-16=15 |
| $\rightarrow 220$ | 12 | REO | 1016 | $016+1 \varnothing 0=1016$ |
| 21 | 14 | STR | 260 |  |
| 22 |  | $\emptyset \emptyset$ | 1 |  |
| 23 | 25 | WIH | 260 |  |
| 24 | 28 | DEI | $\emptyset$ |  |
| L-25 | 21 | JIO | 220 |  |
| 226 | $2 \emptyset$ | JMP | 210 | - |

The upper indexing loop is very similar to that in example C1. In the lower loop the index register is set to 15 first. In the first run REO $1016-1 \varnothing \varnothing \emptyset+15=$ REO 31 acts upon output 31. The instruction DEI (28) decrements the value of the IR by 1 , as a result of which output 30 is switched off in the next run, etc. As soon as the IR $=\emptyset$, no further return is effected, the indexing loop is left.

You will have noticed that in case of the version with relay outputs, a pause is made after 4 relays respectively. This is due to the fact that the 4 unoccupied addresses $2 \emptyset . .23$ are processed too.
(C3) Even calculating is possible with the little PCAØ
The 64 counter registers of the PCAØ can be used in many ways. In connection with the instructions STR and SCR we have only dealt with the codes $\emptyset 0$ and $16,17,18$ of the second line. As evident from the following table, 32 functions are available.

With the codes $\emptyset 1 \ldots 15$ the area from $2 \emptyset 48$ to 65535 can be reached.
With the codes 19...26 8- to 20 -bit digital values can be loaded into the counter register or transferred to elements.

Finally, with the codes $27 \ldots 3 \emptyset$ arithmetic functions can be performed.
Code 31 serves to transmit values from index registers to a counter register or from one counter register to another.

|  | Mnemonic code | Numerical code | Operand | Explanations |
| :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & 1 . \\ & \text { line } \end{aligned}$ | $\text { STR } / \text { SCR }$ | $14 / 15$ | 256... 319 | Address of the register |
| 2. ine |  | $\emptyset \emptyset$ <br> $\emptyset 1$ <br> $\emptyset 2$ <br> $\emptyset 3$ <br> 94 <br> $\emptyset 5$ <br> $\emptyset 6$ <br> 97 <br> 98 <br> $\emptyset 9$ <br> 10 <br> 11 <br> 12 <br> 13 <br> 14 <br> 15 | $\begin{gathered} x x x x \\ \vdots \\ \vdots \\ \vdots \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ \vdots \\ 1 \\ 1 \\ 1 \\ 1 \\ x x x x \end{gathered}$ |  |
|  |  | $\begin{aligned} & 16 \\ & 17 \\ & 18 \end{aligned}$ | Highest addr. \}of 8 subsequent elements | $\begin{array}{rllllr} 2 & \times & 4 & \text { bit } & B C D & \times \\ 2 & \times & 4 & \text { bit } & B C D & \times \\ 2 & \times & 4 & \text { bit } & \text { BCD } & \times \\ 10 \end{array}$ |
|  |  | $\begin{aligned} & 19 \\ & 20 \\ & 21 \\ & 22 \\ & 23 \\ & 24 \\ & 25 \\ & 26 \end{aligned}$ | Highest addr. of the sequence of elements | Read instr. for $5 \times 4$ bit BCD Output instr. for $5 \times 4$ bit BCD Output instr. for 8 bit binary Output instr. for 12 bit binary Output instr. for 16 bit binary Read instr. for 8 bit binary Read instr. for 12 bit binary Read instr. for 16 bit binary |
|  |  | $\begin{aligned} & 27 \\ & 28 \\ & 29 \\ & 30 \end{aligned}$ | $\begin{aligned} & x x x \\ & x x x \\ & x x x \\ & x x x \end{aligned}$ |  |
|  |  | 31 | iii | iii $=\emptyset$ : value of index register is loaded into counter <br> iii $=256 . .319$ : Value of corresponding T/C is loaded into counter |

Problem:
In order to introduce the arithmetical possibilities of the PCAD, we will play a bit with figures in this example. Processing of numerical values plays an important role in counting problems or when analog values have to be processed.

The following functions are assigned to the inputs of our simulation unit:
Iø : Addition +
I1 : Subtraction -
12 : Multiplication $x$
I3 : Division :
I7 : Storing of the 1st value via BCD-preselection switch
I6 : Storing of the 2nd value via BCD-preselection switch
I5 : Triggering of the arithmetic operation
I8... 15 : BCD-preselection switch
C26ø : Register for first value
C27ø : Register for second value
C266 : Display register for DTC
024 (56) : Acknowledgement for storing of 1st value (I7)
025 (57) : Acknowledgement for storing of 2nd value (I6)
026 (58) : Acknowledgement for performing the operation
The individual values stored with I7 and I6 have to be indicated on the operand display. This applies also to the result of the arithmetic operation preselected with I0...I3.

Example: $87-25=62$
. Turn on $\mathrm{I} 1 \longrightarrow$ subtraction

- Set 87 via BCD-switch
- Upon switching on and off I7 the 1st value is loaded into counter C260
. 87 is indicated on the operand display, at the same time LED 24 lights up acknowledging the first step
- Set 25 via BCD-switch
. Upon switching on and off I6 the second value is loaded into counter C27ø
. 25 is indicated on the operand display, at the same time LED 25 lights up acknowledging the second step
- When I5 is on: operation is being performed
- The result of 62 appears on the operand display, at the same time LED 26 lights up acknowledging the third step
. Upon switching off I5 the outputs $24,25,26$ are reset, with I7 a new input can be effected
- If the subtraction had a negative result (2nd figure $>$ 1st figure), 9999 appears on the display
- If in a division a value is divided by $\emptyset, 8888$ is displayed. In both cases, one benefits from the fact that the CPU sets the ACCU $=\emptyset$ in these special cases

Solution:
We use a flow-chart, which enables us to follow the individual steps of the process.

Rough flow-chart


Program:
In order to display the numerical values continuously, we execute DTC 266 in the parallel circulating program PPめ, while the flow-chart containing PP1 is assigned to address 238.

| $\begin{array}{r} \text { ADDR } \\ 23 \emptyset \\ 231 \end{array}$ | NC 29 | MNC PAS $\emptyset \emptyset$ | OPRD 1 238 | $\left\{\begin{array}{l} \text { Assignment of PP1 from } \\ \text { address } 238 \text { on } \end{array}\right.$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 232 | $\emptyset \square$ | NOP | $\emptyset$ |  |  |
| 233 | $\emptyset \emptyset$ | NOP | $\emptyset$ |  |  |
| -234 | 31 | DTC | 266 | 3 PPゆ display C266 |  |
| -235 | 20 | JMP | 234 |  |  |
| - |  |  | 7 | PP1 <br> $\overline{\text { Acknowledgement of step } 1}$ $B C D$ value to $C 260$ |  |
| $\bigcirc 238$ | 26 | WIL | 7 |  |  |
| 39 | 11 | SEO | 24 |  |  |
| 40 | 15 | SCR | 260 |  |  |
| 41 |  | 16 | 15 |  | Storing the 1st value |
| 42 | 15 | SCR | 266 | Copy C260 to C266 for display | (2-digit) |
| 43 |  | 31 | 260) |  |  |
| 44 | 25 | WIH | 7 |  |  |
| 246 | 26 | WIL | 6 | Acknowledgement of step 2 BCD value to C27ø |  |
| 47 | 11 | SE0 | 25 |  |  |
| 48 | 15 | SCR | $27 \emptyset$ |  | Storing the 2nd value (2-digit) |
| 49 |  | 16 | 15 |  |  |
| $5 \emptyset$ | 15 | SCR |  | Copy C27Ø to C266 for display |  |
| 51 |  | 31 | 270) |  |  |
| 52 | 25 | WIH | 6 |  |  |
| 254 | 26 | WIL | 5 | Acknowledgement of step 3$\qquad$ Jump for addition |  |
| 55 | 11 | SEO | 26 |  |  |
| 56 | $\emptyset 1$ | STH | $\emptyset$ |  |  |
| 57 | 21 | JIO | 265 |  |  |
| 58 | $\emptyset 1$ | STH | 1 | Jump for subtraction | $\begin{aligned} & \text { Performing } \\ & \text { operation } \end{aligned}$ |
| 59 | 21 | JIO | $27 \varnothing$ |  |  |
| $6 \emptyset$ | $\emptyset 1$ | STH | 2 | Jump for multiplication |  |
| 61 | 21 | JIO | 280 |  |  |
| 62 | $\emptyset 1$ | STH | 3 |  |  |
| 63 | 21 | JIO | 285 | $\longrightarrow$ Jump for division |  |
| -264 | $2 \emptyset$ | JMP | 295 |  |  |



Manual interrogation or loading of a counter register: If you want to check the contents of a counter register or modify these manually at any time, proceed as follows:
E.g. display C27ø

M Press key "MAN" $\emptyset .5 \mathrm{~s}$
(A) $327 \varnothing$ (that is, $30 \emptyset \emptyset$ is added to the counter address 270)
$\longrightarrow$ Counter contents are displayed

[^3]E.g. enter the value 1234 into C266
(M) Press key "MAN" 0.5 s
(A) 3266 (address + 3øøø)
[E] $\emptyset 1234$ (values < $1 \varnothing \emptyset \varnothing \emptyset$ must be preceded by $\emptyset$ )
$+$

If contact I7 is open, the value introduced in the above program is displayed with DTC, too, in the RUN-mode.

It is hard to believe how many possibilities the little PCAØ offers!
(C4) If long jumps are to be programmed
A program consisting of $2 \not 047$ steps is a long program for a PCAØ. But the user memory has a capacity of $\emptyset . .4 \emptyset 95$ steps.

If long jumps with end addresses in the second half of the user memory, i.e. from $2 \emptyset 48$ to $4 \emptyset 95$ are to be programmed, these jump instructions must be entered using two lines.

This applies to all jump instructions JMP, JIO, JIZ and JMS.
a) Jump instructions with operands 1 to $2 \emptyset 47$

Example: Conditional jump with end address 1845
either $\quad$ JIO (21) $1845 \rightarrow$ one line as usual
or

b) Jump instructions with operands $2 \emptyset 48$ to $4 \emptyset 95$

Example: Jump to the subroutine $328 \emptyset$
For programming:

| 501 502 | JMS (23) | $\emptyset$ 328 | \|E| |
| :---: | :---: | :---: | :---: |
| 501 | JMS (23) | $\emptyset$ | $\square$ |
| 502 | $\emptyset 1$ | 1232 | C |
| 502 | EE | $328 \emptyset$ | + |

The value $\emptyset 11232$ residing in the user memory corresponds to the jump address 3280. 01 stands for the multiple of 2048 and 1232 is the remainder $(1 \times 2 \emptyset 48+1232=328 \emptyset)$.

With key [C the actual jump address is displayed, with the code containing the character EE (applies to the programming unit PCA2.PØ5).

In case of a jump instruction with the operand $\emptyset$ the second line is automatically read for the end address. Therefore, a jump to the address $\emptyset$ always consists of two lines:

c) The following is an example using a blinker in a subroutine. The subroutine starts at address $35 \emptyset \emptyset$.

Main program Subroutine

(C5) The instruction PAS 18 serves to stop the parallel programs, if
they are no longer required

PAS 18 Limitation of the assigned parallel programs
All SAIA-PCs allow assignment of up to 16 parallel programs and running them in parallel. Up to now it has been necessary to. reassign a PP to a dummy program loop if it was no longer required. However, no time could thus be saved during the execution of the remaining PPs.

In case of the PCAD, however, it is now possible to limit a maximum number of active PPS with the PAS 18 -instruction.

After the PP-assignment with PAS $\emptyset . .$. max. 15 a limited number of PPs can be determined in the user program at any place and as many times as desired.

1st line $\quad$| PAS | 18 |
| ---: | ---: |
| 2nd line |  |
| $\emptyset \emptyset$ | $x$ |$; \quad x$ in the range $\emptyset \ldots 15$

Parallel program
Assigned PP
assignment

*) If a higher number than the PPs originally assigned is entered with PAS 18, no malfunction is caused. The inactive PPs 9...11, however, will require processing time in the system program.
(C6) With the instruction "Check-sum" the reliability of your PCAØ is considerably increased

PAS 30 and PAS 31... 34 "Check Sum" of the system and user program
The "Check Sum" serves to establish the check sum of the memory contents of system programs (PAS 30) or of user programs or texts (PAS 31...34). Thus, it can be ensured that the contents of the memories checked have not been changed.

After execution of the instruction:

- $\mathrm{ACCU}=1$ If the comparison is correct
- $\operatorname{ACCU}=\emptyset \quad$ If the reference value does not comply with the check sum.

The instructions PAS 30... 34 are always executed irrespective of the ACCU state.

If a change in contents has occurred, the user can take the measures which seem necessary to him: triggering an alarm, resetting the watchdog etc.

| PAS | $3 \emptyset$ |
| :---: | :---: |
| $\emptyset \emptyset$ | $\emptyset$ |

; Check sum of the system program

| PAS | $31 \ldots 34$ |
| ---: | :--- |
| $x x$ | $x x x x$ |

; 2nd line is always $\emptyset \emptyset \emptyset$
; Check sum of the user program, 31... 34 being introduced corresponding to the program sections 1 K...4K.
; Reference value
The appropriate reference value for the user program is obtained by executing the respective PAS-instruction in the operating mode STEP. The PCAØ displays this check sum on the PCA seconds. In the operating mode PROG, the corresponding reference value can then be introduced in the 2nd line of the PAS 31... 34 instruction.

Attention: Execution of these instructions takes quite a long time:

$$
\begin{array}{ll}
\text { PAS } 30 & =28.0 \mathrm{~ms} \\
\text { PAS } 31 \ldots 34 & =8.3 \mathrm{~ms}
\end{array}
$$

Therefore, use "Check Sum" only if the sequence to be controlled allows it: e.g. when switching on the PC, at the end of an operation cycle, etc.

It is recommended not to introduce this instruction into the user program until it has been completely developed and tested. Each program alteration, irrespective of whether the program was extended or reduced, results in an alteration of the "Check Sum".

Example: A 2 K -user program is to be executed upon switching on


Proceed:

- After entering and checking the
program, select operating mode STEP
- Type in ADR $23+$
$\rightarrow$ the reference value for 2.K (PAS

32) appears for approx. $2 \emptyset \mathrm{sec}$.

- Type ADR 24 (+) for input of
reference value in mode PRG
- Same procedure for PAS 31


## saia'pc

ANLAGE:

| Schritt | Mnemocode | Zahlencode | Operand | Kommentar |
| :---: | :---: | :---: | :---: | :---: |
| $\emptyset$ |  |  |  |  |
| 1 |  |  |  |  |
| 2 |  |  |  |  |
| 3 |  |  |  |  |
| 4 |  |  |  |  |
| 5 |  |  |  |  |
| 6 |  |  |  |  |
| 7 |  |  |  |  |
| 8 |  |  |  |  |
| 9 |  | - |  |  |
| $\emptyset$ |  |  |  |  |
| 1 |  |  |  |  |
| 2 |  |  |  |  |
| 3 |  |  |  |  |
| 4 |  |  |  |  |
| 5 |  |  |  |  |
| 6 |  |  |  |  |
| 7 |  |  |  |  |
| 8 |  |  |  |  |
| 9 |  |  |  |  |
| $\emptyset$ |  |  |  |  |
| 1 |  |  |  |  |
| 2 |  |  |  |  |
| 3 |  |  |  |  |
| 4 |  |  |  |  |
| 5 |  |  |  |  |
| 6 |  |  |  |  |
| 7 |  |  |  |  |
| 8 |  |  |  |  |
| 9 |  |  |  |  |
| $\emptyset$ |  |  |  |  |
| 1 |  |  |  |  |
| 2 |  |  |  |  |
| 3 |  |  |  |  |
| 4 |  |  |  |  |
| 5 |  |  |  |  |
| 6 |  |  |  |  |
| 7 |  |  |  |  |
| 8 |  |  |  |  |
| 9 |  |  |  |  |


| NOP | $\emptyset \emptyset$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| STH | 01 |  |  | X |
| STL | ø2 |  |  | X |
| ANH | Q 3 |  |  | X |
| ANL | ®4 |  |  | X |
| ORH | 05 |  |  | $x$ |
| ORL | 06 |  |  | X |
| XOR | $\emptyset 7$ |  |  | X |
| NEG | 08 |  |  | X |
| DYN | 09 |  |  | X |
| OUT | 1ø |  |  |  |
| SEO | 11 | 1 |  |  |
| REO | 12 | 1 |  |  |
| C00 | 13 | 1 |  |  |
| STR | 14 | 1 |  |  |
| SCR | 15 | 1 |  |  |
| SE I | 16 |  | X | X |
| INC | 17 | 1 |  |  |
| DEC | 18 | 1 |  |  |
| SEA | 19 |  | X | $x$ |
| JMP | 2ø |  | X | $x$ |
| JI0 | 21 |  | X | X |
| JIZ | 22 |  | $x$ | X |
| JMS | 23 |  | X | X |
| RET | 24 |  | X | X |
| WIH | 25 |  | X | X |
| WIL | 26 |  | X | X |
| IN I | 27 |  |  | X |
| DEI | 28 |  |  | $x$ |
| PAS | 29 |  |  |  |
| DOP | 30 | $\emptyset$ |  |  |
| DTC | 31 | 1 |  |  |
| O O O O O E O 들 | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & \stackrel{0}{c} \\ & 0 \\ & \stackrel{\rightharpoonup}{\omega} \\ & \stackrel{N}{N} \end{aligned}$ |  |  | + <br> 0 <br> 0 <br> 0 <br> 4 <br> 4 <br> -1 <br> 0 <br> 0 <br> 0 <br> 0 <br> 0 <br> 3 <br> 3 <br>  |



## ANLAGE：

| Schritt． | Mnemo－ code | Zanlen code | Operand | Kommentar |  |  |  |  |  |  | Schritt |  |  |  |  | Oper |  |  | Kommentar |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Qb |  |  |  |  |  |  |  |  |  |  | 51 |  |  |  |  |  |  |  |  |  |
| 1 |  |  |  |  |  |  |  |  |  |  | 1 |  |  |  |  |  |  |  |  |  |
| 2 |  |  |  |  |  |  |  |  |  |  | 2 |  |  |  |  |  |  |  |  |  |
| 3 |  |  |  |  |  |  |  |  |  |  | 3 |  |  |  |  |  |  |  |  |  |
| 4 |  |  |  |  |  |  |  |  |  |  | 4 |  |  |  |  |  |  |  |  |  |
| 5 |  |  |  |  |  |  |  |  |  |  | 5 |  |  |  |  |  |  |  |  |  |
| 6 |  |  |  |  |  |  |  |  |  |  | 6 |  |  |  |  |  |  |  |  |  |
| 7 |  |  |  |  |  |  |  |  |  |  | 7 |  |  |  |  |  |  |  |  |  |
| 8 |  |  |  |  |  |  |  |  |  |  | 8 |  |  |  |  |  |  |  |  |  |
| 9 |  |  |  |  |  |  |  |  |  |  | 9 |  |  |  |  |  |  |  |  |  |
| 10 |  |  |  |  |  |  |  |  |  |  | 60 |  |  |  |  |  |  |  |  |  |
| 1 |  |  |  |  |  |  |  |  |  |  | 1 |  |  |  |  |  |  |  |  |  |
| 2 |  |  |  |  |  |  |  |  |  |  | 2 |  |  |  |  |  |  |  |  |  |
| 3 |  |  |  |  |  |  |  |  |  |  | 3 |  |  |  |  |  |  |  |  |  |
| 4 |  |  |  |  |  |  |  |  |  |  | 4 |  |  |  |  |  |  |  |  |  |
| 5 |  |  |  |  |  |  |  |  |  |  | 5 |  |  |  |  |  |  |  |  |  |
| 6 |  |  |  |  |  |  |  |  |  |  | 6 |  |  |  |  |  |  |  |  |  |
| 7 |  |  |  |  |  |  |  |  |  |  | 7 |  |  |  |  |  |  |  |  |  |
| 8 |  |  |  |  |  |  |  |  |  |  | 8 |  |  |  |  |  |  |  |  |  |
| 9 |  |  |  |  |  |  |  |  |  |  | 9 |  |  |  |  |  |  |  |  |  |
| $2 \emptyset$ |  |  |  |  |  |  |  |  |  |  | 76 |  |  |  |  |  |  |  |  |  |
| 1 |  |  |  |  |  |  |  |  |  |  | 1 |  |  |  |  |  |  |  |  |  |
| 2 |  |  |  |  |  |  |  |  |  |  | 2 |  |  |  |  |  |  |  |  |  |
| 3 |  |  |  |  |  |  |  |  |  |  | 3 |  |  |  |  |  |  |  |  |  |
| 4 |  |  |  |  |  |  |  |  |  |  | 4 |  |  |  |  |  |  |  |  |  |
| 5 |  |  |  |  |  |  |  |  |  |  | 5 |  |  |  |  |  |  |  |  |  |
| 6 |  |  |  |  |  |  |  |  |  |  | 6 |  |  |  |  |  |  |  |  |  |
| 7 |  |  |  |  |  |  |  |  |  |  | 7 |  |  |  |  |  |  |  |  |  |
| 8 |  |  |  |  |  |  |  |  |  |  | 8 |  |  |  |  |  |  |  |  |  |
| 9 |  |  |  |  |  |  |  |  |  |  | 9 |  |  |  |  |  |  |  |  |  |
| 30 |  |  |  |  |  |  |  |  |  |  | 80 |  |  |  |  |  |  |  |  |  |
| 1 |  |  |  |  |  |  |  |  |  |  | 1 |  |  |  |  |  |  |  |  |  |
| 2 |  |  |  |  |  |  |  |  |  |  | 2 |  |  |  |  |  |  |  |  |  |
| 3 |  |  |  |  |  |  |  |  |  |  | 3 |  |  |  |  |  |  |  |  |  |
| 4 |  |  |  |  |  |  |  |  |  |  | 4 |  |  |  |  |  |  |  |  |  |
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| 6 |  |  |  |  |  |  |  |  |  |  | 6 |  |  |  |  |  |  |  |  |  |
| 7 |  |  |  |  |  |  |  |  |  |  | 7 |  |  |  |  |  |  |  |  |  |
| 8 |  |  |  |  |  |  |  |  |  |  | 8 |  |  |  |  |  |  |  |  |  |
| 9 |  |  |  |  |  |  |  |  |  |  | 9 |  |  |  |  |  |  |  |  |  |
| 40 |  |  |  |  |  |  |  |  |  |  | $9 \varnothing$ |  |  |  |  |  |  |  |  |  |
| 1 |  |  |  |  |  |  |  |  |  |  | 1 |  |  |  |  |  |  |  |  |  |
| 2 |  |  |  |  |  |  |  |  |  |  | 2 |  |  |  |  |  |  |  |  |  |
| 3 |  |  |  |  |  |  |  |  |  |  | 3 |  |  |  |  |  |  |  |  |  |
| 4 |  |  |  |  |  |  |  |  |  |  | 4 |  |  |  |  |  |  |  |  |  |
| 5 |  |  |  |  |  |  |  |  |  |  | 5 |  |  |  |  |  |  |  |  |  |
| 6 |  |  |  |  |  |  |  |  |  |  | 6 |  |  |  |  |  |  |  |  |  |
| 7 |  |  |  |  |  |  |  |  |  |  | 7 |  |  |  |  |  |  |  |  |  |
| 8 |  |  |  |  |  |  |  |  |  |  | 8 |  |  |  |  |  |  |  |  |  |
| 9 |  |  |  |  |  |  |  |  |  |  | 9 |  |  |  |  |  |  |  |  |  |
|  | Logik | $k-B e f e h$ | hle |  |  | ktion | s－Be | fehl |  |  | rung－B | Befe |  |  | ilfs | s－Be | ehl |  | Eingänge Ausqänge | ø－255 |
|  | ｜n｜m | 蔵 | $\infty$ $\hat{\otimes}$ $\infty$ <br> 8   | 98 | $\Rightarrow$ | $\cdots$ | $\pm \xrightarrow{7}$ | $\stackrel{6}{\sim}=$ | － | N | $\cdots$ | $\cdots$ | N | $\stackrel{\sim}{\sim}$ | $\stackrel{\sim}{\sim}{ }^{\infty}$ | －${ }^{\circ}$ | m | $\cdots$ | Zeitglieder <br> Zähler <br> Merker | 256－287 |
| 号京 | 它云云交 |  |  | 2 | 号 |  |  | 式出运 |  | $\sum_{i}^{\circ}$ | 윽 극 | $\sum_{n}^{n} \stackrel{4}{x}$ | $\stackrel{1}{3}$ | 3 | z |  | $\stackrel{\square}{0}$ | 氝 | Haftspeicher | 765－999 |


$\bullet$

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[^0]:    *) By inserting the jumper "NV", all flags and registers for timers and counters are made non-volatile.

[^1]:    *) The values in brackets do not have to be entered, but they are indicated.
    **) For the small PCAØ.M12. . enter output 24, for the big PCAØ.M14.. enter output $4 \varnothing$.

[^2]:    *) These values result from the standard clock rate of $1 / 1 \varnothing \mathrm{~s}$. It can be changed to $1 / 1 \emptyset \emptyset$ s by removing the jumper "1/1ø".

[^3]:    *) As the operand can be at most 2ø47, enter the following (according to the table in example C3):
    instead of $\varnothing 99999 \longrightarrow \emptyset 41807(4 \times 2048+1807)$
    instead of $\emptyset \varnothing 8888 \longrightarrow \emptyset 4696(4 \times 2 \emptyset 48+696)$

